DETAILED ACTION

This office action is in response to amendments, remarks and request for reconsideration filed on August 17, 2011. The remarks and amendments filed have been entered and made of record. Claims 29-30, 32, 34-36, 38-43 45-48 and 50-58 are pending.

Continued Examination Under 37 CFR 1.114

A request for continued examination under 37 CFR 1.114, including the fee set forth in 37 CFR 1.17(e), was filed in this application after final rejection. Since this application is eligible for continued examination under 37 CFR 1.114, and the fee set forth in 37 CFR 1.17(e) has been timely paid, the finality of the previous Office action has been withdrawn pursuant to 37 CFR 1.114. Applicant's submission filed on August 17, 2011 has been entered

Response to Arguments

Applicant's arguments filed on August 17, 2011 have been fully considered but they are not persuasive because of the following reasons and newly applied art Hearn et al. (WO 03/003242):

Applicant's arguments filed on October 2, 2009 have been fully considered but they are not persuasive because of the following reasons:

Regarding Claims 29 applicants argued, that system of cited prior arts (CPA) [Thibadeau U. S. Patent 7,036,020 and Hearn et al. (WO 03/003242).] does not disclose or teach *"selectively blocking access to operating system data and instead of writing the changes to the operating system files in the security partition as targeted by the requests, writes the changes to a location different than the security partition"*.

The system of cited prior art teaches a system and method for promoting security method in computer system that involves partitioning portion of storage device to form security partition and limiting access to portion of storage device by operating system of computer. In that system a portion of storage device is partitioned to form a security partition, which has an authority record and data set associated with the authority record. An access to security partition of storage device is limited by the installed operating system of computer. Thibadeau :(Fig.1-4, col.4 line 37 to col.6 line 16). While Hearn teaches and describe a blocking unit (Fig.1 item 25 and Fig.2) configured for host CPU to impose and continuously maintain the requisite level of data access to security partition for users effecting the data access in accordance with the particular data access profile, regardless of subsequent operations of the CPU (Hearn: Fig.1 Item 35, Fig.2 and page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As a result, cited prior art does implement and teach a system that relates to securing access in a computer system. Therefore, the examiner asserts that cited prior art does teach or suggest the subject matter broadly recited in independent Claims and in subsequent dependent Claims. Accordingly, rejections for claims 29-30, 32, 34-36, 38-43 45-48 and 50-58 are respectfully maintained.

Claim Rejections - 35 USC § 103

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all

obviousness rejections set forth in this Office action:

> (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in
> section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are
> such that the subject matter as a whole would have been obvious at the time the invention was made to a person
> having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the
> manner in which the invention was made.

Claims 29-30, 32, 34-36, 38-43 45-48 and 50-58 are rejected under 35 U.S.C. 103(a) as

being unpatentable over Thibadeau U. S. Patent 7,036,020 and further in view of Hearn et al.

(WO 03/003242).

2.   Regarding Claim 29, Thibadeau teaches and describes a security system for securing

access to an operating system of a computer having at least a host central processing unit (CPU),

a memory used by the host CPU to load programs from the operating system in order to operate

the computer, a storage device for storing data to be used by the computer, and a chain of

components connecting the CPU to the storage device, the security system comprising:

a security partition formed in the storage device, the operating system of the computer being

stored in the security partition; and a security device comprising a hardware processor or

controller for intercepting communications and selectively blocking access to operating system

data between the host CPU and the security partition; wherein the security device is deployed

along the chain of components that connect the host CPU to the storage device; wherein the

security device's processor or controller is distinct from the host CPU; and wherein during

operation of the operating system the security device is arranged to: intercept requests to write

changes to operating system files in the security partition; in response to the requests, instead of

write the changes to the operating system files in the security partition as targeted by the

requests, write the changes to a location different than the security partition; and cause normal

operation of the operating system to continue without writing the changes to the operating

system files in the security partition (Fig.1-4, col.4 line 37 to col.6 line 16).

Although the system disclosed by Thibadeau shows all the features of the claimed

limitation, but Thibadeau does not specifically disclose security device processor controller for

selectively blocking access to operating system data between CPU and the security partition.

In an analogous art, Hearn, on the other hand, discloses a security device comprising a

hardware processor or controller for intercepting communications and selectively blocking

access to operating system data between the host CPU and the security partition, wherein the

security device is deployed along the chain of components that connect the host CPU to the

storage device, wherein the security device's processor or controller is distinct from the host CPU

[Fig.1 Item 35, Fig.2 and page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8].

Therefore, It would have been obvious to one ordinary skilled in the art at the time of

invention to combine the teachings of Thibadeau and Hearn, because Hearn method security

device would not only promote further security structure and control in the system of Thibadeau

during receiving data from user or devices but will also provide safeguards and imposing

requisite level of data access to security partition for users from unauthorized access and misuse

(page 35 line 28 to page 36 line 8).

3.   Regarding Claim 42, Thibadeau teaches and describes a method for securing access to an

operating system of a computer, comprising: forming a security partition in the storage device;

storing the operating system of the computer in the security partition; loading operating system

data from the operating system into a random access memory; using one or more host central

processing units (CPUs)to execute programs in the operating system based on the operating

system data loaded into the random access memory; intercepting communications and selectively

blocking access to operating system data between the host CPUs and the security partition at a

security device deployed along the chain of components connecting the host CPUs to the storage

device, wherein the security device operates independent of the host CPU; and intercepting, at

the security device, requests to write changes to operating system files in the security partition; in

response to the requests, instead of writing the changes to the operating system files in the

security partition as targeted by the requests, the security device writing the changes to a location

different than the security partition; and

        the security device causing normal operation of the operating system to continue without

writing the changes to the operating system files in the security partition (Fig.1-4, col.4 line 37 to

col.6 line 16).

Although the system disclosed by Thibadeau shows all the features of the claimed limitation, but

Thibadeau does not specifically disclose security device processor controller for selectively

blocking access to operating system data between CPU and the security partition.

        In an analogous art, Hearn, on the other hand, discloses a security device comprising a

hardware processor or controller for intercepting communications and selectively blocking

access to operating system data between the host CPU and the security partition, wherein the

security device is deployed along the chain of components that connect the host CPU to the

storage device, wherein the security device's processor or controller is distinct from the host CPU

[Fig.1 Item 35, Fig.2 and page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8].

Therefore, It would have been obvious to one ordinary skilled in the art at the time of

invention to combine the teachings of Thibadeau and Hearn, because Hearn method security

device would not only promote further security structure and control in the system of Thibadeau

during receiving data from user or devices but will also provide safeguards and imposing

requisite level of data access to security partition for users (page 35 line 28 to page 36 line 8).

4.   Claims 30, 32, 34-36, 38-41, 43, 45-48 and 50-54 are rejected applied as above rejecting

Claims 29, and 42. Furthermore, the system of Thibadeau and Hearn teaches and describes a

system and method for securing access to an operating system of a computer, wherein:

As per Claim 30, each user of the computer has an associated access profile, each access

profile comprising information indicative of the level of access to portions of the storage device

permitted by a user, and the security device controlling access to the storage device by a user in

accordance with the access profile associated with the user (col. 6 line 55 to col.8 line 35 and

Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 32, said security device is independent and separately configurable of said

host CPU (col.4 line 37 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to

page to page 17 line 8).

As per Claim 34, the location different than the security partition is at least a portion of a flash ROM (Fig.1-4, col.4 line 37 to col.5 line 50 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 35, the location different than the security partition is at least a portion of an invisible partition formed in the storage device invisible partition formed in the storage device (col.5 line 15 to col.6 line 16).

As per Claim 36, further comprising authentication means for authenticating a user of the computer and associating the user with a prescribed access profile, said security device controlling subsequent access to the security partition in accordance with the access profile associated with the user (col. 6 line 55 to col.8 line 35 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 38, said blocking means is configured to block all access by the host CPU to the storage device before initialisation of the security system, and to selectively permit access immediately after said initialisation in accordance with a respective access profile (col. 6 line 55 to col.8 line 35 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 39, said authentication means enables a software boot of the computer to be effected only after correct authentication of a user, and said security system permits normal loading of the operating system during the start up sequence of the computer following said software boot (col.6 line 55 to col.8 line 35 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 40, said security device is a security device physically deployed between an interface adapter and the storage device within a data access channel of the chain of components

connecting the host CPU and the storage device (col.4 line 37 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 41, said security device is integrated in a bridging circuit within the chain of components connecting the host CPU and the storage device or within the storage device (Fig.1-4, and col.4 line 37 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 43, further comprising associating each user with an access profile comprising information indicative of the level of access to portions of the storage device permitted by a user; and for each user, selectively blocking access between the host CPU and the security partition in accordance with the access profile defined for the user (col.5 line 25 to col.6 line 16, and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 45, further comprising authenticating a user of the computer, and associating the user with an access profile after successful user authentication (col.5 line 15 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 46, said selective blocking comprises controlling access between the host CPU and the security partition independently of the host CPU (col.4 line 37 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 47, said selective blocking comprises totally blocking access to the storage device by the host CPU during initialisation of the computer, and intercepting all said access immediately after said initialisation and before loading of the operating system of the computer (col.6 line 55 to col.8 line 35, and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8)).

As per Claim 48, including performing a software boot of the computer only after correct

authentication of the user, and allowing normal loading of the operating system during the start

up sequence of the computer after said software boot (col. 6 line 55 to col.8 line 35 and Hearn:

page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 50, to location different than the security partition is at least a portion of

flash ROM in the security device (Fig.1-4, col.4 line 37 to col.5 line 50 and Hearn: page 15 line

2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 51, the location different than the security partition is at least a portion of

an invisible partition formed in the storage device (Fig.1-4, col.4 line 37 to col.6 line 16 and

Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 52.including unalterably storing computer programs for effecting said

controlling access in a location separate from the memory and not addressable by the host CPU

(Fig.1-4, and col.4 line 37 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to

page to page 17 line 8).

As per Claim 53, the security device is a dedicated hardware device comprising a

dedicated CPU for processing the intercepted communication and, based on the intercepted

communications, determining whether to block data access between the host CPU and the

security partition (col.4 line 45 to line 65, and col.9 line 13 to line 22, and Hearn: page 15 line 2

to line 8, page 5 line 24 to page to page 17 line 8)).

As per Claim 54, the security device is integrated into a bridging circuit comprising logic

for processing the intercepted communications and, based on the intercepted communications,

determining whether to block data access between the host CPU and the security partition (col.4

line 45 to line 65, and col.9 line 13 to line 22, and Hearn: page 15 line 2 to line 8, page 5 line 24

to page to page 17 line 8).

As per Claim 55, the security device is arranged to divert to the scratch location write

operations to the security partition requested by a first user, but permit performance of write

operations to the security partition requested by a second user (col.5 line 25 to col.6 line 16 and

Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page 17 line 8).

As per Claim 56, the security device is further configured to erase the changes written to

the scratch location without the changes having been written to the operating system files

targeted by the requests (col.5 line 25 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5

line 24 to page to page 17 line 8).

As per Claim 57, further comprising diverting to the scratch location write operations to

the security partition requested by a first user, but permitting performance of write operations to

the security partition requested by a second user (col.5 line 25 to col.6 line 16 and Hearn: page

15 line 2 to line 8, page 5 line 24 to page to page 17 line 8)..

As per Claim 58, further comprising erasing the changes written to the scratch location

without the changes having been written to the operating system files targeted by the requests

(col.5 line 25 to col.6 line 16 and Hearn: page 15 line 2 to line 8, page 5 line 24 to page to page

17 line 8).

Conclusion

Any inquiry concerning this communication or earlier communications from the examiner should be directed to SYED ZIA whose telephone number is (571)272-3798. The examiner can normally be reached on 9:00 to 5:00.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Nathan J. Flynn can be reached on571-272-1915. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

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October 9, 2011
/Syed   Zia/
Primary Examiner, Art Unit 2431